

# Explicit isogenies in quadratic time in any characteristic

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# Summary

- 1 Explicit isogenies
- 2 Couveignes' algorithm
- 3 An  $\ell$ -adic Couveignes algorithm (special case)
- 4 An  $\ell$ -adic Couveignes algorithm

# Refresher on isogenies

## Isogenies

- $E$  and  $E'$  two elliptic curves, an isogeny is a surjective morphism  $\phi : E \rightarrow E'$  such that  $\phi(0_E) = 0_{E'}$ . Isogenies are
  - group morphisms (with finite kernel),
  - algebraic maps (of finite degree).
- If the isogeny has degree  $\ell$ , we call it an  $\ell$ -isogeny,
- We say that  $E$  and  $E'$  are  $\ell$ -isogenous if there exist an  $\ell$ -isogeny between them.
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**Example:** Let  $\ell$  be an integer, the *scalar multiplication map*

$$\begin{aligned} [\ell] : E &\rightarrow E \\ P &\mapsto \ell \cdot P \end{aligned}$$

is an isogeny of degree  $\ell^2$ .

# Explicit isogenies

## Explicit isogenies

Separable isogenies of degree  $\ell$  are represented by rational maps in Weierstrass form

$$(x, y) \mapsto \left( \frac{n(x)}{d(x)}, cy \left( \frac{n(x)}{d(x)} \right)' \right)$$

with  $n(x)$  a polynomial of degree  $\ell$  and  $d(x)$  a polynomial of degree  $\ell - 1$ .

## Vélu's formulas

From the knowledge of its kernel we can explicitly write down the isogeny:

- Let  $G = \ker \phi$ , then  $d(X) = \prod_{P \in G \setminus \{0_E\}} (X - x(P))$ ;
- Vélu's formulas: from  $d(X)$  we deduce  $n(X)/d(X)$ ;

# Motivation

## Explicit isogeny computation problem

Let  $\mathbb{F}_q$  be a finite field of characteristic  $p$ . Given  $E, E'$  two  $r$ -isogenous elliptic curves defined over  $\mathbb{F}_q$ , compute an  $r$ -isogeny  $\phi : E \rightarrow E'$ .

Applications:

- Schoof-Elkies-Atkin point counting algorithm,
- ECC cryptanalysis: [Gaudry, Hess, Smart '02],
- Hash functions: [Charles, Goren, Lauter '07],
- Trapdoors: [Teske '06],
- Post quantum cryptography: [Rotostev, Stolbunov '06], [De Feo, Jao, Plût '11].

# Previous work

Let  $p$  be the characteristic of  $\mathbb{F}_q$ .

- [Elkies '92/'98], [Bostan, Morain, Salvy, Schost '08] use  $\tilde{O}(r)$  operations in  $\mathbb{F}_q$ , work only for  $r < 2p$ . Specific to the SEA case.
  - [Couveignes '94] any characteristic,  $\tilde{O}(r^3 p^{O(1)})$  operations.
  - [Lercier '97] only  $p = 2$ .
  - [Couveignes '96], [De Feo '10] any characteristic,  $\tilde{O}(r^2 p^{O(1)})$  operations.
  - [Lercier, Sirvent '08], [Lairez, Vaccon '16] works for every  $p$  using  $\tilde{O}(r^2)$  operations in  $\mathbb{F}_q$ . Specific to the SEA case.
- We want to modify Couveignes' algorithm to obtain an algorithm with complexity  $\tilde{O}(r^2)$  but with no exponential dependency in  $\log(p)$ .

# Torsion points of elliptic curves

## Torsion points

Let  $E$  be an elliptic curve defined over a finite field  $\mathbb{F}_q$ , and let  $m > 0$

$$E[m] = \{P \in E(\bar{\mathbb{F}}_q), mP = 0_E\}$$

For *ordinary* elliptic curves

$$E[\ell^k] \simeq \mathbb{Z}/\ell^k\mathbb{Z} \times \mathbb{Z}/\ell^k\mathbb{Z} \quad \text{with } \ell \neq p$$

$$E[p^k] \simeq \mathbb{Z}/p^k\mathbb{Z}$$



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## Couveignes' algorithm (compute an $r$ -isogeny $\phi : E \rightarrow E'$ )

Compute  $\phi$  by interpolation over  $E[p^k]$ :

- Compute generators  $P, P'$  of  $E[p^k], E'[p^k]$ ;
- Interpolate  $\phi$ , assuming it maps  $uP \mapsto uP'$  for all  $u \in \mathbb{Z}/p^k\mathbb{Z}$ ;
- Test whether  $\phi$  is an isogeny.

In case it is not, replace  $P'$  with a multiple  $aP'$  and start again.

# Couveignes algorithm (1996)

**Input:**  $E, E'$  two  $r$ -isogenous curves on  $\mathbb{F}_{p^n}$

**Output:**  $\phi : E \rightarrow E'$  of degree  $r$

- 1 Select the least  $k$  such that  $p^k > 4r$ ;
- 2 Compute generators  $P$  of  $E[p^k]$  and  $P'$  of  $E'[p^k]$ ;
- 3 Compute  $T = \prod (X - x(uP))$  with  $1 \leq u \leq \frac{p^k-1}{2}$ ;
- 4 For each  $a \in (\mathbb{Z}/p^k\mathbb{Z})^\times$ :  $O(r)$ 
  - 1 Compute the interpolation polynomial  $L$  such that  $L(x(uP)) = x(a u P')$  for all  $u \in \mathbb{Z}/p^k\mathbb{Z}$ ;  $\tilde{O}(rp^{O(1)})$
  - 2 Use a rational reconstruction algorithm to compute a rational fraction  $F = L \bmod T$  of degrees  $(r, r-1)$ ;  $\tilde{O}(r)$
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## Goal

Replace  $E[p^k]$  by  $E[\ell^k]$  for a small prime  $\ell \neq p$ .

## An $\ell$ -adic Couveignes' algorithm?

Our goal is to work with  $E[\ell^k] \simeq (\mathbb{Z}/\ell^k\mathbb{Z})^2$  instead of  $E[p^k]$  to remove the polynomial dependency in  $p$ .

- $E[p^k] = \langle P \rangle \simeq (\mathbb{Z}/p^k\mathbb{Z})$  with  $p^k \approx r$
- $E[\ell^k] = \langle P, Q \rangle \simeq (\mathbb{Z}/\ell^k\mathbb{Z}) \times (\mathbb{Z}/\ell^k\mathbb{Z})$  with  $\ell^{2k} \approx r$

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## $p$ -adic

Let  $P \in E$  and  $P' \in E'$

$$P \mapsto aP' \quad a \in (\mathbb{Z}/p^k\mathbb{Z})^*$$

$\Rightarrow O(r)$  possibilities.

## $\ell$ -adic

Let  $P, Q \in E$  and  $P', Q' \in E'$

$$\begin{pmatrix} P \\ Q \end{pmatrix} \mapsto \begin{pmatrix} a & b \\ c & d \end{pmatrix} \begin{pmatrix} P' \\ Q' \end{pmatrix}$$

with  $\begin{pmatrix} a & b \\ c & d \end{pmatrix} \in \text{GL}_2(\mathbb{Z}/\ell^k\mathbb{Z})$  invertible.

$\Rightarrow O(r^2)$  possibilities.

# Frobenius vs isogenies

## Definition (Frobenius Endomorphism)

$E$  an ordinary elliptic curve defined over  $\mathbb{F}_q$ . The function

$$\pi : (x, y) \mapsto (x^q, y^q)$$

is called Frobenius endomorphism. It satisfies a quadratic equation

$$\pi^2 - t_\pi \pi + q = 0.$$

We are only working with rational isogenies  $\psi : E \rightarrow E'$ , i.e.

$$\pi_{E'} \circ \psi = \psi \circ \pi_E.$$

Subgroup of size  $\ell$



$\ell$ -isogeny

Subgroup of size  $\ell$   $\Leftrightarrow$   $\ell$ -isogeny

Subgroup of size  $\ell$  stable by  $\pi$   $\Leftrightarrow$  Rational  $\ell$ -isogeny



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Assume that  $\pi$  splits modulo  $\ell$ : i.e. its minimal polynomial factors as

$$(\pi - \lambda)(\pi - \mu) \quad \text{with} \quad \lambda \neq \mu \pmod{\ell}$$

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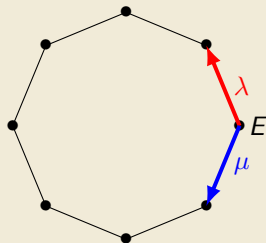
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Two eigenspaces in  $E[\ell]$   $\Rightarrow$  Two rational  $\ell$ -isogenies  
 $\ker(\pi - \lambda), \ker(\pi - \mu)$  of direction  $\lambda, \mu$

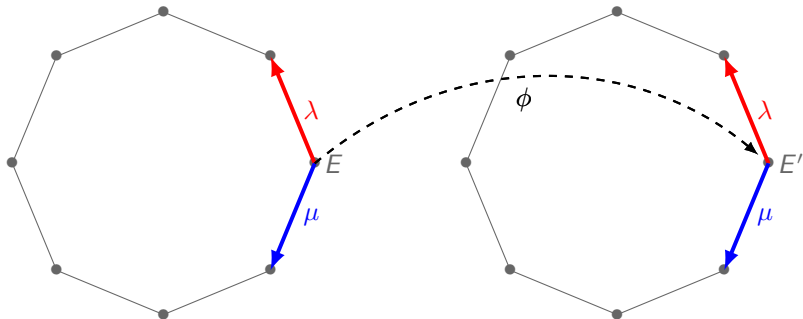
## Isogeny graph



## Fact

Let  $\phi$  be an  $r$ -isogeny with  $\ell \nmid r$ , then  $\phi$  preserves the kernels of the  $\ell$ -isogenies of direction  $\lambda, \mu$ .

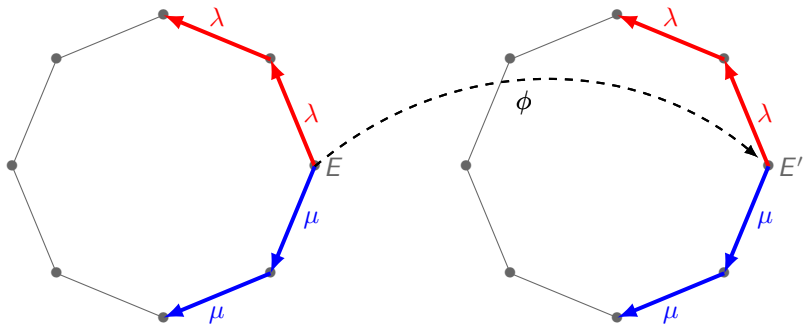
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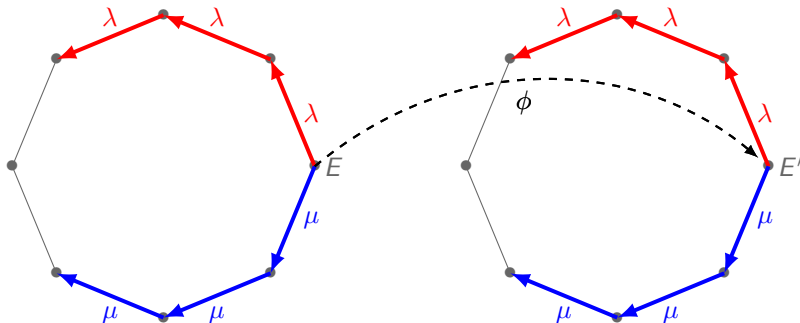
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To interpolate  $\phi$  over  $E[\ell^k]$ , we want to compute two cyclic  $\ell^k$ -subgroups of direction  $\lambda, \mu$ .



- We call  $E[\ell^k]_\lambda \oplus E[\ell^k]_\mu$  a *horizontal decomposition*;
- SEA literature calls this an *isogeny cycle* [Couveignes, Morain '94].

## Towards an $\ell$ -adic Couveignes' algorithm ( $\pi$ splits modulo $\ell$ )

**Input:**  $E, E'$  two  $r$ -isogenous curves on  $\mathbb{F}_q$

**Output:**  $\phi : E \rightarrow E'$  of degree  $r$

**Fact:**  $\phi$  maps  $E[\ell^k]_\lambda \rightarrow E'[\ell^k]_\lambda$  and  $E[\ell^k]_\mu \rightarrow E'[\ell^k]_\mu$ .

- 1 Select the least  $k$  such that  $\ell^{2k} > 4r$ .
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and  $\langle P', Q' \rangle = E'[\ell^k]_\lambda \oplus E'[\ell^k]_\mu$
- 3 For each **invertible diagonal** matrix  $\begin{pmatrix} a & 0 \\ 0 & b \end{pmatrix}$  in  $(\mathbb{Z}/\ell^k\mathbb{Z})^{2 \times 2}$ :
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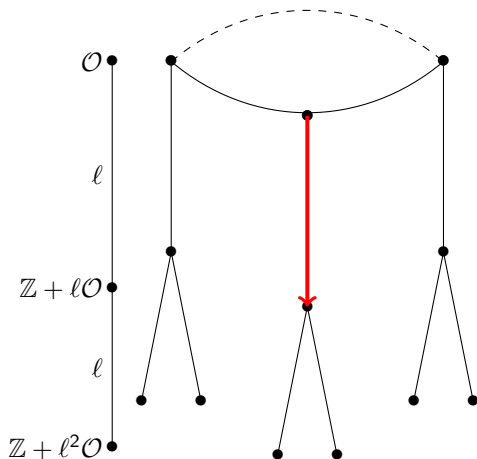
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# The general case

Denote by  $\mathcal{O}$  (resp.  $\mathcal{O}'$ ) the endomorphism ring of  $E$  (resp.  $E'$ )



## Lemma (Kohel 1996)

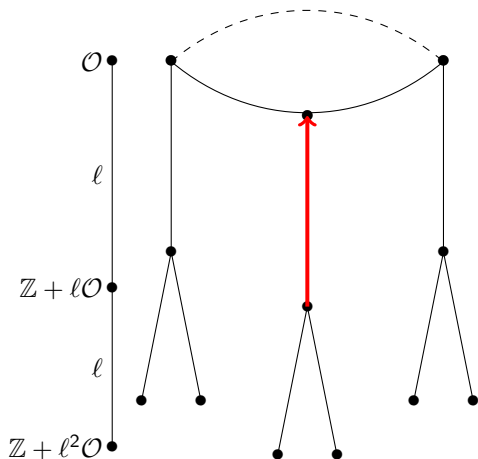
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 Then we say that  $\psi$  is

- 1 **descending** if  $\ell = [\mathcal{O} : \mathcal{O}']$
- 2 **ascending** if  $\ell = [\mathcal{O}' : \mathcal{O}]$ ,
- 3 **horizontal** if  $\mathcal{O} = \mathcal{O}'$ .



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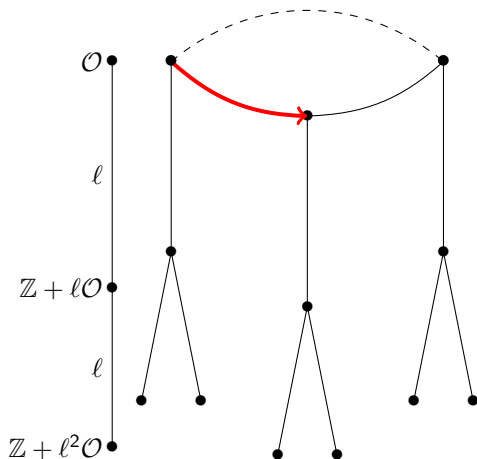
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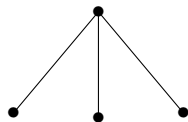


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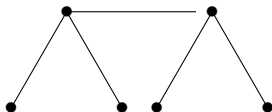
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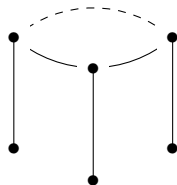
# A guide to volcano types



Inert prime  $\ell$



Ramified prime  $\ell$



Split prime  $\ell$

Figure: The three shapes of volcanoes of 2-isogenies

In the rest of this talk we consider only volcanoes with cyclic crater (Elkies case).

## Elkies prime

We say that  $\ell$  is an *Elkies prime* if the characteristic polynomial of  $\pi$  factors over  $\mathbb{Z}_\ell$  as

$$\pi^2 - t_\pi \pi + q = (\pi - \lambda)(\pi - \mu), \quad \text{with } \lambda \neq \mu,$$

where  $h = v_\ell(\lambda - \mu)$  can be  $\geq 1$ .

Since for any  $P \in E[\ell^h]$  we have:

$$\pi(P) = \lambda P = \mu P$$

we cannot immediately distinguish kernels of isogenies with direction  $\lambda$  from those with direction  $\mu$ .

Thus we have to work with  $E[\ell^{h+1}]$  to get  $P \in E[\ell^{h+1}]_\lambda$  such that:

$$\pi(P) = \lambda P \neq \mu P$$

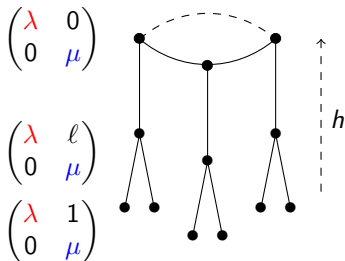
We must therefore now assume that  $k \geq h + 1$ .

## Proposition (De Feo, H., Plût, Schost)

In the Elkies case the action of the Frobenius endomorphism  $\pi$  on  $E[\ell^{h+1}]$  is conjugate, over  $\mathbb{Z}_\ell$ , to a unique matrix

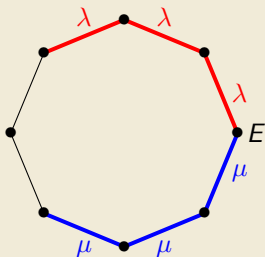
$$\begin{pmatrix} \lambda & a \\ 0 & \mu \end{pmatrix},$$

with  $a \in \{1, \ell, \dots, \ell^{h-1}, 0\}$ , and  $a = 0$  iff  $E$  lies on the crater.



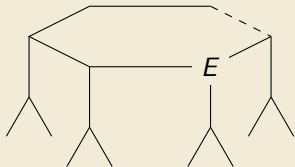
We assume that  $E$  lies on the crater (we can reduce to this case easily).

### Volcano with height $h = 0$



$\ker(\pi - \mu \mid E[\ell^k])$  is a cyclic group of size  $\ell^k$ .

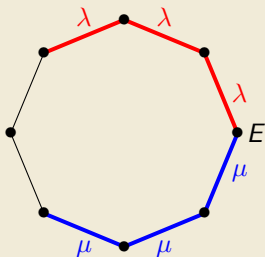
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**Problem:** if  $h \geq 1$  then  $\ker(\pi - \mu \mid E[\ell^k]) \simeq (\mathbb{Z}/\ell^k) \times (\mathbb{Z}/\ell^h)$  is not cyclic, and contains  $\ell^h$  cyclic subgroups of order  $\ell^k$ .

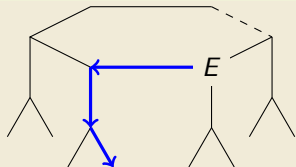
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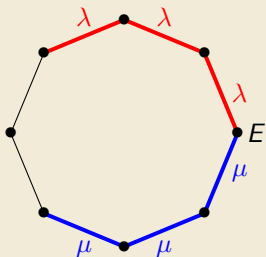
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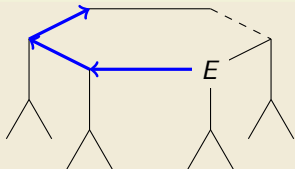
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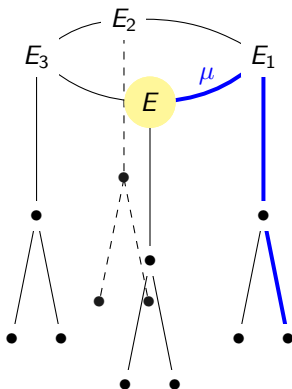
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# Walking on the crater

**Goal:** compute the kernel of horizontal  $\ell^k$ -isogenies (codomain curves lying on the crater).

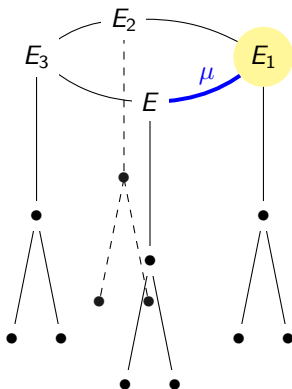
We proceed step by step by walking along the volcano crater according to the direction  $\mu$ .



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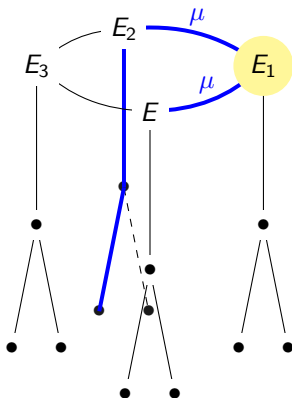
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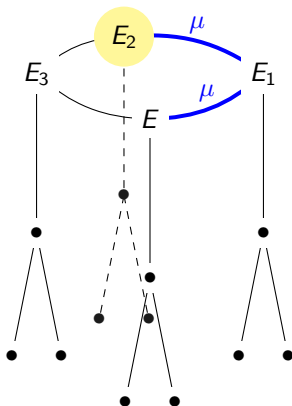
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# Walking on the crater

**Goal:** compute the kernel of horizontal  $\ell^k$ -isogenies (codomain curves lying on the crater).

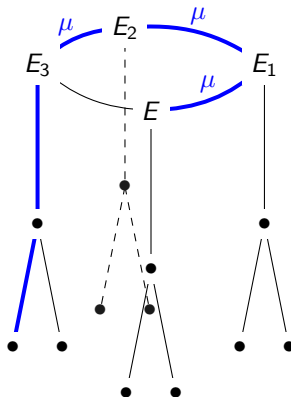
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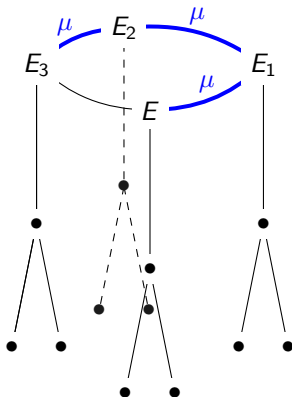
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# Technical details

## Computing in towers of field extensions

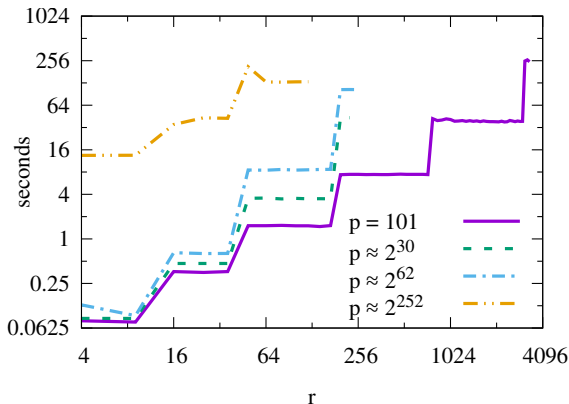
- Torsion points are not defined in  $\mathbb{F}_q$ , in general.
- We work in  $\ell$ -adic extensions of  $\mathbb{F}_q$  using constructions from [De Feo, Doliskani, Schost '13], [Doliskani, Schost '15] where in particular we have a fast computation of the Frobenius.

## Finding an Elkies prime $\ell$

- The complexity depends polynomially on the auxiliary prime  $\ell$ .
- Ideally we would like to work with  $\ell = 2$ .
- In practice half of all  $\ell$  are expected to be Elkies primes.
- In theory we can only prove  $\ell \leq O(\log(q))$  for almost all  $q$  and curves  $E, E'$  (see [Shparlinski, Sutherland '14]).

# Experiments

The algorithm has been implemented on SageMath v7.1 for the case of  $\ell = 2$ , the code is available on GitHub:  
[https://github.com/Hugounenq-Cyril/Two\\_curves\\_on\\_a\\_volcano](https://github.com/Hugounenq-Cyril/Two_curves_on_a_volcano)





# Conclusion

## Contribution

- New tools for navigating isogeny volcanoes.
- A faster variant of Couveignes' algorithm.

## Future work

- Compare implementation to other algorithms (esp. Lercier-Sirvent).
- Give an analogous algorithm for Atkin primes.
- Analyze our techniques to navigate the volcano in other settings: point counting, computation of endomorphism rings, Hilbert class polynomials, modular polynomials.

# A GUIDE TO VOLCANO TYPES

